

# Module Learning with Errors

## with General Distributions

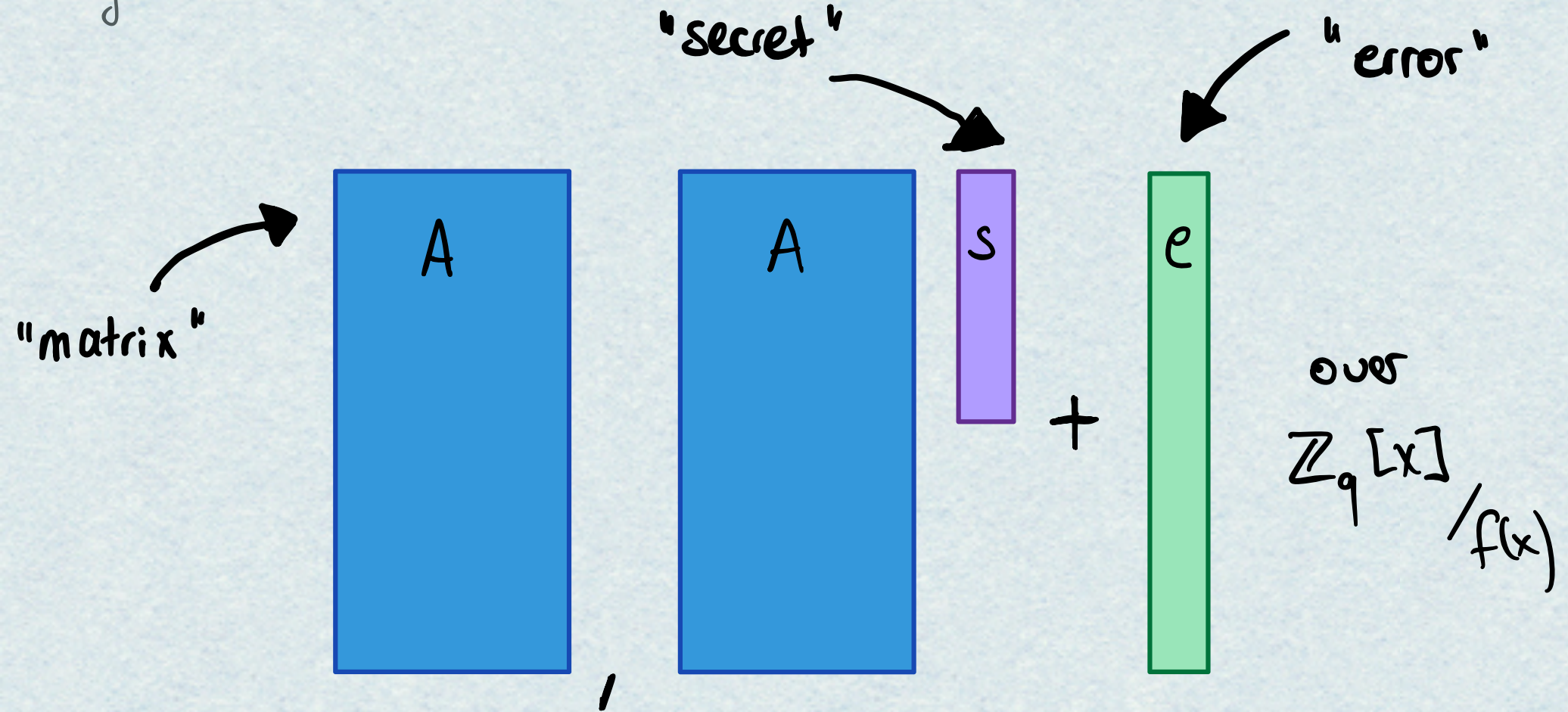
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# Module Learning With Errors (M-LWE)

Langlois, Stehlé DCC'15

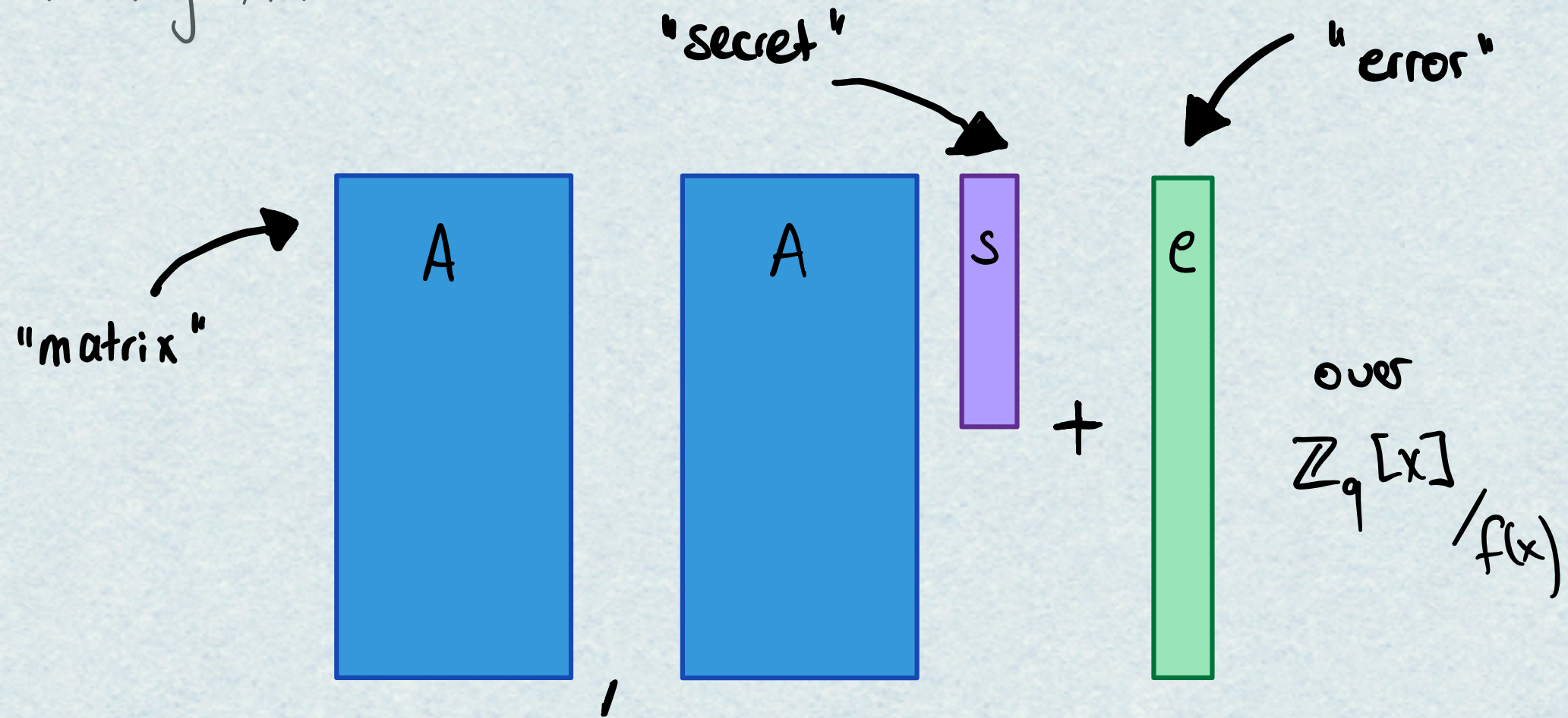


Search: Find **s** (or **e**)

Decision: Distinguish from  $(A, \text{uniform})$

# Module Learning With Errors (M-LWE)

Roux-Langlois, Stehlé DCC'18



Choices:

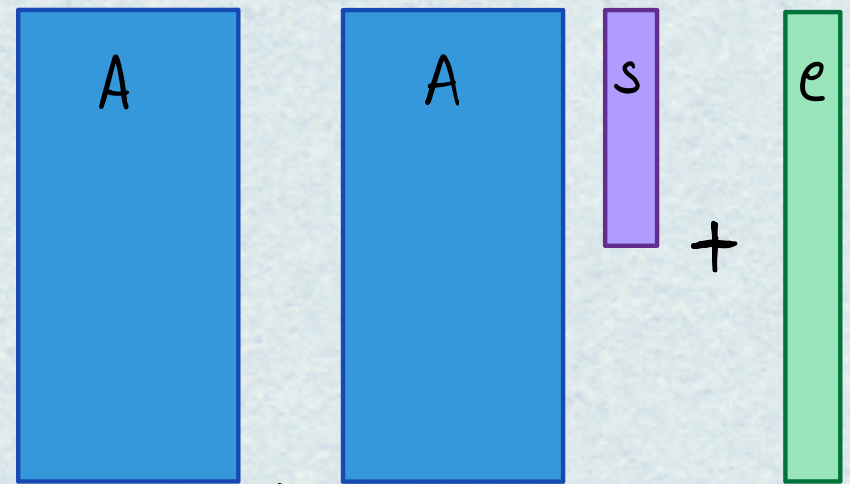
- \* polynomial  $f(x)$
- \* distribution of secret  $s$
- \* distribution of error  $e$
- \* distribution of matrix  $A$

} flexible usage  
BUT  
non-trivial  
security analysis

# Standard Module Learning With Errors

Short Independent Vectors Problem  
in ANY module lattice

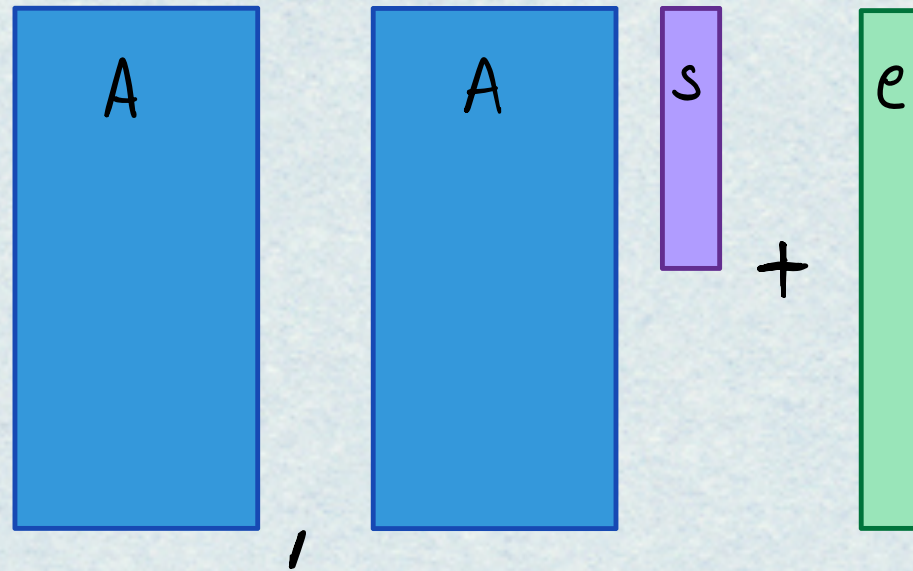
↓  
quantum: Roux-Langlois, Stehlé ICC'15  
classical: Boudgoust, Joux, Roux-Langlois, Wen Asiacrypt'20



- \* cyclotomic polynomial  $f(x)$
- \*  $s$  uniform over  $\mathbb{Z}_q[x]/f(x)$
- \*  $e$  discrete Gaussian
- \*  $A$  uniform over  $\mathbb{Z}_q[x]/f(x)$

Variants of  
M-LWE

# Module Learning with Errors Variants



\* distribution of  $s$ : any as long as enough min-entropy

Brakerski, Dötting TCC'20

Boudgoust, Joux, Roux-Langlois, Wen Indocrypt'22

Lin, Wang, Zhuang, Wang TCS'24

Question: Can we show something similar for the distribution of  $e$ ?

## Our Results Part 1:

Search  $\mu$ -LWE with noise distribution  $\mathcal{D}$   
is as hard as standard  $\mu$ -LWE  
as long as:

- ① samples of  $\mathcal{D}$  are short
- ②  $\mathcal{D}$  has sufficient min-entropy
- ③ not too many samples are provided

Application: sparse ternary error distribution

# Module Learning With Errors in Hermite Normal Form

Applebaum, Cash, Peikert, Sahai Crypto'09

$s$  follows the same  
distribution as  $e$

- both are short
- relevant for efficiency

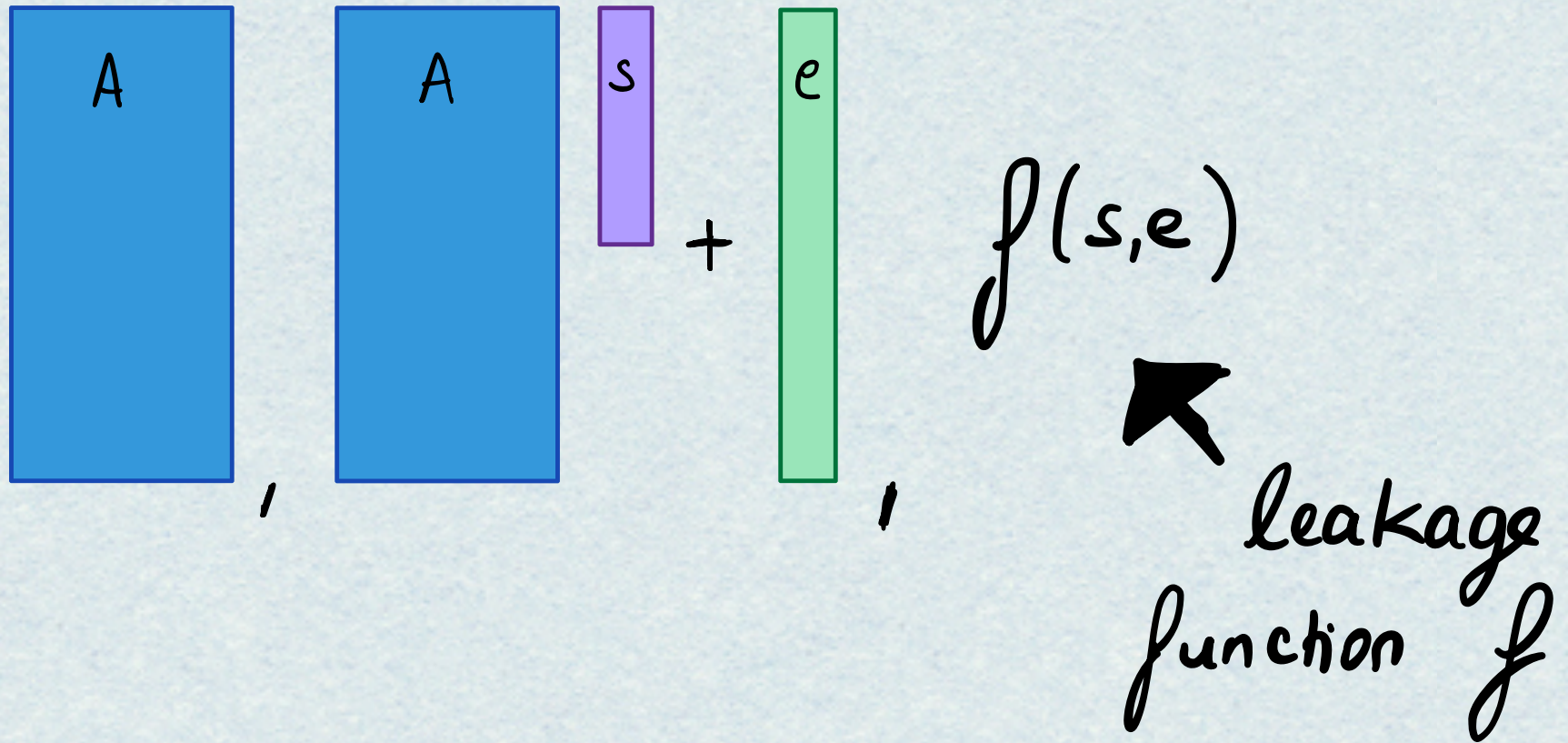
## Our Results Part 2:

Search  $M$ -LWE with secret-noise distr.  $\mathcal{D}$   
is as hard as standard  $M$ -LWE  
as long as:

- ① samples of  $\mathcal{D}$  are short
- ②  $\mathcal{D}$  has sufficient min-entropy
- ③ not too many, but still enough samples are provided

We address a subtlety if the coefficients of  $(\frac{s}{e}) \in \mathcal{D}$  are correlated.

# Module Learning With Errors *Leaky* Variants



- \* within reductions
- \* leverage constructions

# Our Results Part 3:

Search  $M$ -LWE with secret-noise distr.  $\mathcal{D}$   
and leakage  $f(s, e)$

is as hard as standard  $M$ -LWE

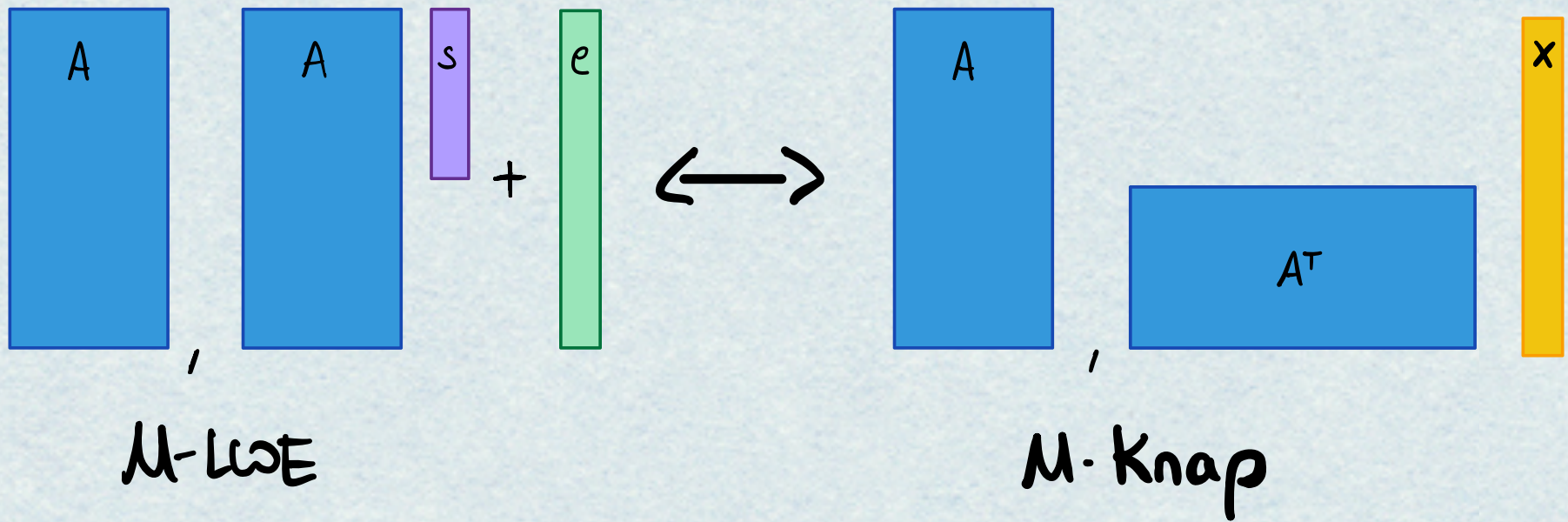
as long as:

- ① samples of  $\mathcal{D}$  are short
- ②  $\mathcal{D}$  has sufficient min-entropy  
conditioned on leakage  $f$
- ③ not too many, but still enough  
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Detailed Analysis for

- \* exact & approximate linear  $f$ 's
- \* some quadratic function
- \* some non-algebraic function

# Proof Strategy for general error distribution



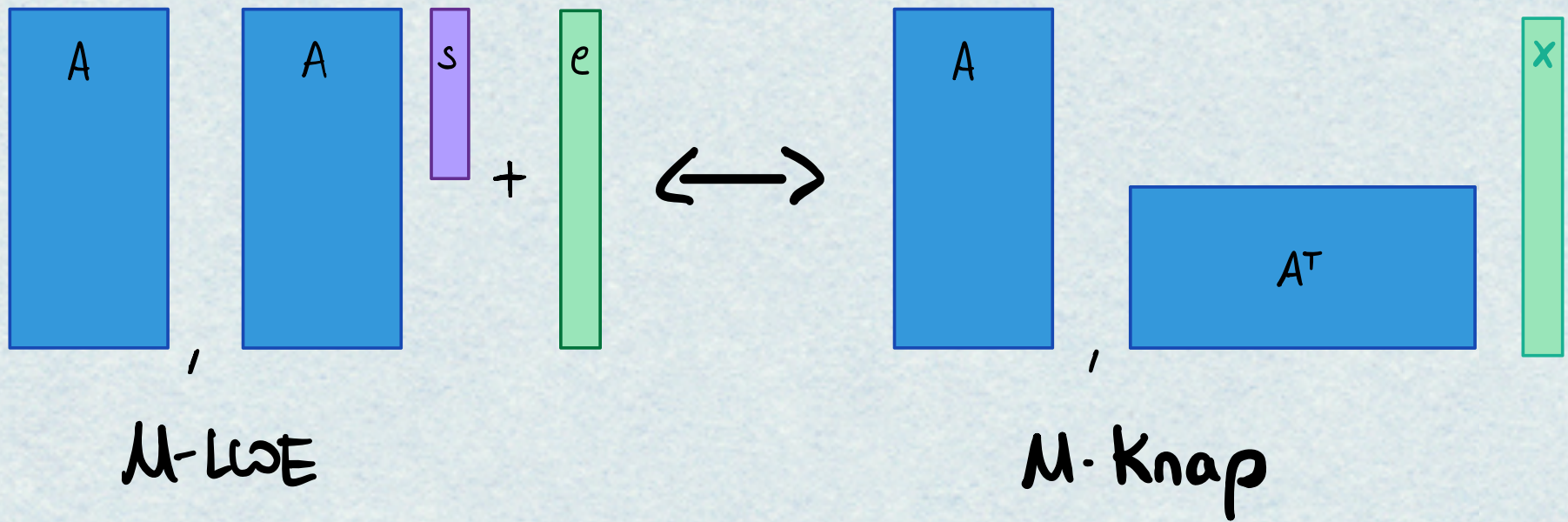
Function family  $f_A(x) = A^T x$

Closely follows proofs for bounded uniform error

LWE: Micciancio & Peikert Crypto'13

M-LWE: Boudgoust, Joux, Roux-Langlois, Wen JoC'23

# Proof Strategy for general error distribution



Function family  $f_A(x) = A^T x$

Second preimage resistant

uninvertible

$(f_A)_A$

one-way

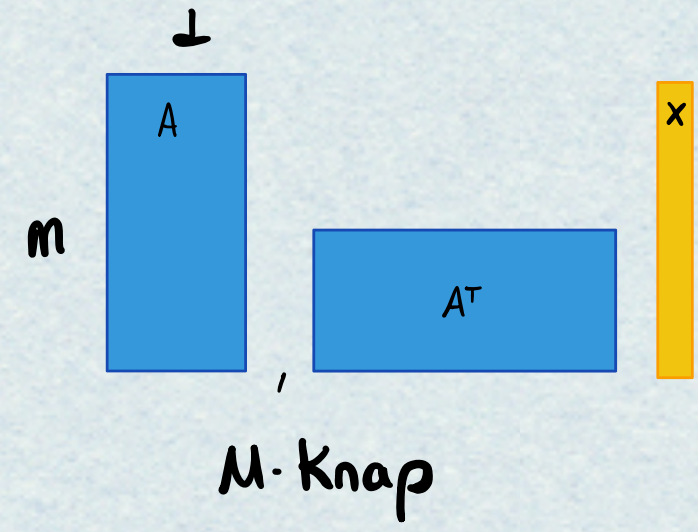
search  $\Leftrightarrow$  M-Knap is hard

# Second-Preimage Resistance of $\mathcal{M}$ -Knap

Given  $(A, x \leftarrow \mathcal{D})$

find  $x' \neq x$  st  $A^T x = A^T x'$

$x' \in \text{Supp}(\mathcal{D}) = S$



$$\sum_{x' \in S \setminus \{x\}} \mathbb{P}_A [A^T(x - x') = 0]$$

$$\leq (|S| - 1) \cdot \frac{\text{Bound}^d}{|R_q|^d}$$

loose

uniformly distributed over the ideal

$$I = \langle x_1 - x'_1, \dots, x_m - x'_m \rangle$$

$$\mathcal{N}(I) \in \text{Bound}$$

## Wrapping-UP:

Search  $M$ -LWE with secret-noise distr.  $\mathcal{D}$   
and leakage  $f(s, e)$

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## Limitations:

- \* practical parameters not covered
- \* only search variant

## Wrapping-up:

# Thanks!

Search  $M$ -LWE with secret-noise distr.  $\mathcal{D}$   
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